

Notes on Nova

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Abstract

Notes taken while reading Nova [1] paper.
Usually while reading papers I take handwritten notes, this document contains some of them re-written to *LaTeX*.
The notes are not complete, don't include all the steps neither all the proofs.

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1 Folding Scheme for Committed Relaxed R1CS

1.1 R1CS modification

Want: merge 2 instances of R1CS with the same matrices into a single one. Each instance has $z_i = (W_i, x_i)$ (public witness, private values resp.).

traditional R1CS Merged instance with $z = z_1 + rz_2$, for rand r . But, since R1CS is not linear \rightarrow can not apply.

eg.

$$\begin{aligned} Az \circ Bz &= A(z_1 + rz_2) \circ B(z_1 + rz_2) \\ &= Az_1 \circ Bz_1 + r(Az_1 \circ Bz_2 + Az_2 \circ Bz_1) + r^2(Az_2 \circ Bz_2) \\ &\neq Cz \end{aligned}$$

\rightarrow introduce error vector $E \in \mathbb{F}^m$, which absorbs the cross-terms generated by folding.

\rightarrow introduce scalar u , which absorbs an extra factor of r in $Cz_1 + r^2Cz_2$ and in $z = (W, x, 1 + r \cdot 1)$.

Relaxed R1CS

$$u = u_1 + ru_2$$

$$E = E_1 + r(Az_1 \circ Bz_2 + Az_2 \circ Bz_1 - u_1Cz_2 - u_2Cz_1) + r^2E_2$$

$$Az \circ Bz = uCz + E, \text{ with } z = (W, x, u)$$

where R1CS set $E = 0$, $u = 1$.

$$\begin{aligned} Az \circ Bz &= Az_1 \circ Bz_1 + r(Az_1 \circ Bz_2 + Az_2 \circ Bz_1) + r^2(Az_2 \circ Bz_2) \\ &= (u_1Cz_1 + E_1) + r(Az_1 \circ Bz_2 + Az_2 \circ Bz_1) + r^2(u_2Cz_2 + E_2) \\ &= u_1Cz_1 + \underbrace{E_1 + r(Az_1 \circ Bz_2 + Az_2 \circ Bz_1) + r^2E_2}_{E} + r^1u_2Cz_2 \\ &= u_1Cz_1 + r^2u_2Cz_2 + E \\ &= (u_1 + ru_2) \cdot C \cdot (z_1 + rz_2) + E \\ &= uCz + E \end{aligned}$$

For R1CS matrices (A, B, C) , the folded witness W is a satisfying witness for the folded instance (E, u, x) .

Problem: not non-trivial, and not zero-knowledge. Solution: use polynomial commitment with hiding, binding, succinctness and additively homomorphic properties.

Committed Relaxed R1CS Instance for a Committed Relaxed R1CS (\bar{E}, u, \bar{W}, x) , satisfied by a witness (E, r_E, W, r_W) such that

$$\bar{E} = Com(E, r_E)$$

$$\bar{W} = Com(W, r_W)$$

$$Az \circ Bz = uCz + E, \text{ where } z = (W, x, u)$$

1.2 Folding protocol

V and P take two *committed relaxed R1CS* instances

$$\varphi_1 = (\bar{E}_1, u_1, \bar{W}_1, x_1)$$

$$\varphi_2 = (\bar{E}_2, u_2, \bar{W}_2, x_2)$$

P additionally takes witnesses to both instances

$$(E_1, r_{E_1}, W_1, r_{W_1})$$

$$(E_2, r_{E_2}, W_2, r_{W_2})$$

Let $Z_1 = (W_1, x_1, u_1)$ and $Z_2 = (W_2, x_2, u_2)$.

1. P send $\bar{T} = Com(T, r_T)$,
 where $T = Az_1 \circ Bz_1 + Az_2 \circ Bz_2 - u_1 Cz_2 - u_2 Cz_2$
 and rand $r_T \in \mathbb{F}$
2. V sample random challenge $r \in \mathbb{F}$
3. V, P output the folded instance (\bar{E}, u, \bar{W}, x)

$$\bar{E} = \bar{E}_1 + r\bar{T} + r^2\bar{E}_2$$

$$u = u_1 + ru_2$$

$$\bar{W} = \bar{W}_1 + r\bar{W}_2$$

$$x = x_1 + rx_2$$

4. P outputs the folded witness (E, r_E, W, r_W)

$$E = E_1 + rT + r^2E_2$$

$$r_E = r_{E_1} + r \cdot r_T + r^2 r_{E_2}$$

$$W = W_1 + rW_2$$

$$r_W = r_{W_1} + r \cdot r_{W_2}$$

P uses a zkSNARK showing that knows the valid witness (E, r_E, W, r_W) for the committed relaxed R1CS without revealing its value. Then, via Fiat-Shamir transform we achieve non-interactivity.

2 IVC proofs

WIP

References

- [1] Abhiram Kothapalli, Srinath Setty, and Ioanna Tzialla. Nova: Recursive zero-knowledge arguments from folding schemes. Cryptology ePrint Archive, Paper 2021/370, 2021. <https://eprint.iacr.org/2021/370>.